

POWER ITERATION

1. Reduction to triangular form

The relevance of eigendecompositions $A = X \Lambda X^{-1}$ to repeated application of the linear operator $A \in \mathbb{C}^{m \times m}$ as in

$$e^{tA} = I + \frac{1}{1!}tA + \frac{1}{2!}t^2A^2 + \dots = X e^{t\Lambda} X^{-1},$$

suggests that algorithms that construct powers of A might reveal eigenvalues. This is indeed the case and leads to a class of algorithms of wide applicability in scientific computation. First, observe that taking condition numbers gives

$$\mu(A) = \mu(X \Lambda X^{-1}) \leq \mu^2(X) \mu(\Lambda) = (|\lambda|_{\max} / |\lambda|_{\min}),$$

where $|\lambda|_{\max}$, $|\lambda|_{\min}$ are the eigenvalues of maximum and minimum absolute values. While these express an intrinsic property of the operator A , the factor $\mu^2(X)$ is associated with the conditioning of a change of coordinates, and a natural question is whether it is possible to avoid any ill-conditioning associated with a basis set X that is close to linear dependence. The answer to this line of inquiry is given by the following result.

SCHUR THEOREM. For any $A \in \mathbb{C}^{m \times m}$ there exists Q unitary and T upper triangular such that $A = QTQ^*$.

Proof. Proceed by induction, starting from an arbitrary eigenvalue λ and eigenvector x . Let $u_1 = x / \|x\|$, the first column vector of a unitary matrix $U = [u_1 \ V]$. Then

$$U^*AU = \begin{bmatrix} u_1^* \\ V^* \end{bmatrix} A [u_1 \ V] = \begin{bmatrix} u_1^* \\ V^* \end{bmatrix} [Au_1 \ AV] = \begin{bmatrix} u_1^* \\ V^* \end{bmatrix} [\lambda u_1 \ AV] = \begin{bmatrix} \lambda_1 & b^* \\ \mathbf{0} & C \end{bmatrix},$$

with $C \in \mathbb{C}^{(m-1) \times (m-1)}$ that by the inductive hypothesis can be written as $C = WSW^*$, with W unitary, S upper triangular. The matrix

$$Q = U \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W \end{bmatrix}$$

is a product of unitary matrices, hence itself unitary. The computation

$$Q^*AQ = \left(U \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W \end{bmatrix} \right)^* AU \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W \end{bmatrix} = \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W^* \end{bmatrix} U^*AU \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W \end{bmatrix} = \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W^* \end{bmatrix} \begin{bmatrix} \lambda_1 & b^* \\ \mathbf{0} & C \end{bmatrix} \begin{bmatrix} 1 & \mathbf{0} \\ \mathbf{0} & W \end{bmatrix} = \begin{bmatrix} \lambda_1 & b^* \\ \mathbf{0} & S \end{bmatrix} = T,$$

then shows that T is indeed triangular. □

The eigenvalues of an upper triangular matrix are simply its diagonal elements, so the Schur factorization is an eigenvalue-revealing factorization.

2. Power iteration for real symmetric matrices

When the operator A expresses some physical phenomenon, the principle of action and reaction implies that $A \in \mathbb{R}^{m \times m}$ is symmetric, $A = A^T$ and has real eigenvalues. Componentwise, symmetry of $A = [a_{ij}]$ implies $a_{ij} = a_{ji}$. Consider $Ax = \lambda x$, and take the adjoint to obtain $x^T A^T = \bar{\lambda} x^T$, or $x^T A = \bar{\lambda} x^T$ since A is symmetric. Form scalar products $x^T Ax = \lambda x^T x$, $x^T A^T x = \bar{\lambda} x^T x$, and subtract to obtain

$$0 = (\lambda - \bar{\lambda})x^T x \Rightarrow \lambda = \bar{\lambda} \Rightarrow \lambda \in \mathbb{R},$$

since $x \neq \mathbf{0}$, an eigenvector.

Example. Consider a linear array of identical mass-springs. The i^{th} point mass obeys the dynamics

$$m\ddot{x}_i = k(x_{i+1} - x_i) - k(x_i - x_{i-1}) = k(x_{i+1} - 2x_i + x_{i-1}),$$

expressed in matrix form as $\ddot{x} = Ax$, with A symmetric.

For a real symmetric matrix the Schur theorem states that

$$\mathbf{A} = \mathbf{A}^T \Rightarrow (\mathbf{Q}\mathbf{T}\mathbf{Q}^T) = \mathbf{Q}\mathbf{T}^T\mathbf{Q}^T \Rightarrow \mathbf{T} = \mathbf{T}^T,$$

and since a symmetric triangular matrix is diagonal, the Schur factorization is also an eigendecomposition, and the eigenvector matrix \mathbf{Q} is a basis, $C(\mathbf{Q}) = \mathbb{R}^m$.

2.1. The power iteration idea

Assume initially that the eigenvalues are distinct and ordered $|\lambda_1| > |\lambda_2| > \dots > |\lambda_m|$. Repeated application of \mathbf{A} on an arbitrary vector $\mathbf{v} = \mathbf{Q}\mathbf{c} \in \mathbb{R}^m = C(\mathbf{Q})$ is expressed as

$$\mathbf{A}^n \mathbf{v} = (\mathbf{Q}\mathbf{\Lambda}\mathbf{Q}^T)^n \mathbf{Q}\mathbf{c} = (\mathbf{Q}\mathbf{\Lambda}\mathbf{Q}^T)(\mathbf{Q}\mathbf{\Lambda}\mathbf{Q}^T) \dots (\mathbf{Q}\mathbf{\Lambda}\mathbf{Q}^T) \mathbf{Q}\mathbf{c} = \mathbf{Q}\mathbf{\Lambda}^n \mathbf{c},$$

a linear combination of the columns of \mathbf{Q} (eigenvectors of \mathbf{A}) with coefficients $\mathbf{\Lambda}^n \mathbf{c} = [\lambda_1^n c_1 \ \lambda_2^n c_2 \ \dots \ \lambda_m^n c_m]^T$.

- For large enough n , $|\lambda_1| > |\lambda_k|$, $k = 2, \dots, n$, leads to a dominant contribution along the direction of eigenvector \mathbf{q}_1

$$\mathbf{A}^n \mathbf{v} = \mathbf{Q}\mathbf{\Lambda}^n \mathbf{c} = \lambda_1^n c_1 \mathbf{q}_1 + \dots + \lambda_m^n c_m \mathbf{q}_m \approx \lambda_1^n c_1 \mathbf{q}_1.$$

This gives a procedure for finding one eigenvector of a matrix, and the Schur theorem proof suggests a recursive algorithm to find all eigenvalues can be defined.

The sequence of normalized eigenvector approximants $\mathbf{v}_n = \mathbf{A}^n \mathbf{v} / \|\mathbf{A}^n \mathbf{v}\|$ is linearly convergent at rate $r = |\lambda_2 / \lambda_1|$.

2.2. Rayleigh quotient

To estimate the eigenvalue revealed by power iteration, formulate the least squares problem

$$\min_c \|\mathbf{A}\mathbf{v} - \mathbf{v}c\|,$$

that seeks the best approximation of one power iteration $\mathbf{A}\mathbf{v}$ as a linear combination of the initial vector \mathbf{v} . Of course, if $\mathbf{v} = \mathbf{q}$ is an eigenvector, then the solution would be $c = \lambda$, the associated eigenvalue. The projector onto $C(\mathbf{v})$ is

$$\mathbf{P} = \frac{\mathbf{v}\mathbf{v}^T}{\mathbf{v}^T \mathbf{v}},$$

that when applied to $\mathbf{A}\mathbf{v}$ gives the equation

$$\mathbf{P}\mathbf{A}\mathbf{v} = \frac{\mathbf{v}\mathbf{v}^T}{\mathbf{v}^T \mathbf{v}} \mathbf{A}\mathbf{v} = \frac{\mathbf{v}^T \mathbf{A}\mathbf{v}}{\mathbf{v}^T \mathbf{v}} \mathbf{v} = c\mathbf{v} \Rightarrow c = \frac{\mathbf{v}^T \mathbf{A}\mathbf{v}}{\mathbf{v}^T \mathbf{v}}.$$

The function $r: \mathbb{R}^m \rightarrow \mathbb{R}$,

$$r(\mathbf{v}) = \frac{\mathbf{v}^T \mathbf{A}\mathbf{v}}{\mathbf{v}^T \mathbf{v}},$$

is known as the Rayleigh quotient which, evaluated for an eigenvector, gives $r(\mathbf{q}) = \lambda$. To determine how well the eigenvalue is approximated, carry out a Taylor series in the vicinity of an eigenvector \mathbf{q}

$$r(\mathbf{v}) = r(\mathbf{q}) + \frac{1}{1!} [\nabla_{\mathbf{v}} r(\mathbf{q})]^T (\mathbf{v} - \mathbf{q}) + \mathcal{O}(\|\mathbf{v} - \mathbf{q}\|^2),$$

where $\nabla_{\mathbf{v}} r$ is the gradient of $r(\mathbf{v})$

$$\nabla_{\mathbf{v}} r = \begin{bmatrix} \frac{\partial r}{\partial v_1} \\ \vdots \\ \frac{\partial r}{\partial v_m} \end{bmatrix}.$$

Compute the gradient through differentiation of the Rayleigh quotient

$$\nabla_{\mathbf{v}} r(\mathbf{v}) = \frac{\nabla_{\mathbf{v}} (\mathbf{v}^T \mathbf{A}\mathbf{v})}{\mathbf{v}^T \mathbf{v}} - \frac{(\mathbf{v}^T \mathbf{A}\mathbf{v})}{(\mathbf{v}^T \mathbf{v})^2} \nabla_{\mathbf{v}} (\mathbf{v}^T \mathbf{v}).$$

Noting that $\nabla_{\mathbf{v}} v_i = \mathbf{e}_i$, the i^{th} column of \mathbf{I} , the gradient of $\mathbf{v}^T \mathbf{v}$ is

$$\nabla_{\mathbf{v}} (\mathbf{v}^T \mathbf{v}) = \nabla_{\mathbf{v}} \sum_{i=1}^m v_i^2 = \sum_{i=1}^m \nabla_{\mathbf{v}} v_i^2 = \sum_{i=1}^m 2v_i \nabla_{\mathbf{v}} v_i = 2 \sum_{i=1}^m v_i \mathbf{e}_i = 2\mathbf{v}.$$

To compute $\nabla_{\mathbf{v}}(\mathbf{v}^T \mathbf{A} \mathbf{v})$, let $\mathbf{u} = \mathbf{A} \mathbf{v}$, and since \mathbf{A} is symmetric $\mathbf{u}^T = \mathbf{v}^T \mathbf{A}^T = \mathbf{v}^T \mathbf{A}$, leading to

$$\nabla_{\mathbf{v}}(\mathbf{v}^T \mathbf{A} \mathbf{v}) = \nabla_{\mathbf{v}}(\mathbf{u}^T \mathbf{v}) = \sum_{i=1}^m \nabla_{\mathbf{v}}(u_i v_i) = \sum_{i=1}^m u_i \nabla_{\mathbf{v}} v_i + \sum_{i=1}^m v_i \nabla_{\mathbf{v}} u_i.$$

Use $u_i = \sum_{j=1}^m a_{ij} v_j$ also expressed as $u_j = \sum_{i=1}^m a_{ji} v_i$ by swapping indices to obtain

$$\nabla_{\mathbf{v}} u_i = \sum_{j=1}^m a_{ij} \nabla_{\mathbf{v}} v_j = \sum_{j=1}^m a_{ij} \mathbf{e}_j$$

and therefore

$$\sum_{i=1}^m v_i \nabla_{\mathbf{v}} u_i = \sum_{i=1}^m v_i \sum_{j=1}^m a_{ij} \mathbf{e}_j = \sum_{j=1}^m \sum_{i=1}^m a_{ij} v_i \mathbf{e}_j = \sum_{j=1}^m \sum_{i=1}^m a_{ij} v_i \mathbf{e}_j.$$

Use symmetry of \mathbf{A} to write

$$\sum_{i=1}^m a_{ij} v_i = \sum_{i=1}^m a_{ji} v_i = u_j,$$

and substitute above to obtain

$$\sum_{i=1}^m v_i \nabla_{\mathbf{v}} u_i = \sum_{j=1}^m u_j \mathbf{e}_j = \mathbf{u} = \mathbf{A} \mathbf{v}.$$

Gathering the above results

$$\nabla_{\mathbf{v}}(\mathbf{v}^T \mathbf{v}) = 2\mathbf{v}, \nabla_{\mathbf{v}}(\mathbf{v}^T \mathbf{A} \mathbf{v}) = 2\mathbf{A} \mathbf{v},$$

gives the following gradient of the Rayleigh quotient

$$\nabla_{\mathbf{v}} r(\mathbf{v}) = \frac{2}{\mathbf{v}^T \mathbf{v}} (\mathbf{A} \mathbf{v} - r(\mathbf{v}) \mathbf{v}).$$

When evaluated at $\mathbf{v} = \mathbf{q}$, obtain $\nabla_{\mathbf{v}} r(\mathbf{q}) = \mathbf{0}$, implying that near an eigenvector the Rayleigh quotient approximation of an eigenvalue is of quadratic accuracy,

$$r(\mathbf{v}) - \lambda = \mathcal{O}(\|\mathbf{v} - \mathbf{q}\|^2).$$

2.3. Refining the power iteration idea

Power iteration furnishes the largest eigenvalue. Further eigenvalues can be found by use of the following properties:

- (λ, \mathbf{q}) eigenpair of $\mathbf{A} \Rightarrow (\lambda - \mu, \mathbf{q})$ eigenpair of $\mathbf{A} - \mu \mathbf{I}$;
- (λ, \mathbf{q}) eigenpair of $\mathbf{A} \Rightarrow (1/\lambda, \mathbf{q})$ eigenpair of \mathbf{A}^{-1} .

If μ is a known initial approximation of the eigenvalue then the inverse power iteration $\mathbf{v}_n = (\mathbf{A} - \mu \mathbf{I})^{-1} \mathbf{v}_{n-1}$, actually implemented as successive solution of linear systems

$$(\mathbf{A} - \mu \mathbf{I}) \mathbf{v}_n = \mathbf{v}_{n-1},$$

leads to a sequence of Rayleigh quotients $r(\mathbf{v}_n)$ that converges quadratically to an eigenvalue close to μ . An important refinement of the idea is to change the shift at each iteration which leads to cubic order of convergence

Algorithm (Rayleigh quotient iteration)

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Given  $\mathbf{v}, \mathbf{A}$ 
 $\mu = \mathbf{v}^T \mathbf{A} \mathbf{v} / \mathbf{v}^T \mathbf{v}$ 
for  $i = 1$  to  $n_{\max}$ 
   $\mathbf{w} = (\mathbf{A} - \mu \mathbf{I}) \setminus \mathbf{v}$  (solve linear system)
   $\mathbf{v} = \mathbf{w} / \|\mathbf{w}\|$ 
   $\lambda = \mathbf{v}^T \mathbf{A} \mathbf{v}$ 
  if  $|\lambda - \mu| < \varepsilon$  exit
   $\mu = \lambda$ 
end
return  $\lambda, \mathbf{v}$ 

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Power iteration can be applied simultaneously to multiple directions at once

Algorithm (Simultaneous iteration)

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Given  $A$   
 $Q = I; \mu = \text{diag}(A)$   
for  $i = 1$  to  $n_{\max}$   
   $V = AQ$  (power iteration applied to multiple directions)  
   $QR = V$  (orthogonalize new directions)  
   $\lambda = \text{diag}(Q^T A Q)$   
  if  $\|\lambda - \mu\| < \varepsilon$  exit  
end  
return  $\lambda, Q$ 
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