## Test 2

Solve the following problems (6 course points each). Present a brief motivation of your method of solution.

1. Consider the variant of Horner's algorithm for evaluation of a polynomial  $p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_0$ that saves intermediate results in  $(p_n, p_{n-1}, \dots, p_0) \in \mathbb{R}^{n+1}$ , with final result returned in  $p_0$ 

Input:  $x, a_n, ..., a_0$   $p_n = a_n$ for i = n - 1 downto 0  $p_i = x \cdot p_{i+1} + a_i$ end return $(p_0, ..., p_n)$ 

Admiting the floating point arithmetic axiom  $x \circledast y = (x \ast y)(1 + \epsilon)$  for all  $x, y \in \mathbb{F}$  (set of real floating point numbers), find an expression for  $\delta p = p(x + \delta x) - p(x)$  with  $\delta x$  denoting the ensemble of floating point errors introduced by every arithmetic operation in Horner's algorithm

- a) Is polynomial evaluation backward stable?
- b) Show that the relative condition number of polynomial evaluation is well approximated by

$$\kappa = \frac{|a_0| + |a_1x| + \dots + |a_nx^n|}{|a_0 + a_1x + \dots + a_nx^n|}$$

- c) Is the evaluation of z(x) = 0 well-conditioned, ill-conditioned or ill-posed?
- d) Define the relative distance between two polynomials p, q with coefficients  $a = (a_0, a_1, ..., a_n), b = (b_0, b_1, ..., b_n)$  as

$$d(p,q) = \max_{i} \left| \frac{a_i - b_i}{a_i} \right|.$$

Show that the inverse condition number of evaluating p is also the minimal distance to an ill-posed polynomial evaluation.

Solution. Assuming that a machine epsilon floating point error  $\epsilon$  is introduced by each operation, at stage i

$$\tilde{p}_i = (x \cdot \tilde{p}_{i+1}(1+\epsilon) + a_i)(1+\epsilon) \Rightarrow \tilde{p}_i = x \cdot \omega^2 \cdot \tilde{p}_{i+1} + \omega \cdot a_i = y \cdot \tilde{p}_{i+1} + b_i$$

with  $y = \omega^2 x$ ,  $b_i = \omega \cdot a_i$  ( $b_n = \omega \cdot a_n$  is also included to force symmetric treatment of polynomial coefficients, and can be considered the floating point result of  $1 \cdot a_n$ ). The effect of floating point errors in Horner's algorithm is to evaluate the polynomial

$$\tilde{q}(y) = b_n y^n + b_{n-1} y^{n-1} + \dots + b_0.$$

Substitute to obtain expression in terms of  $x, a_0, ..., a_n$ 

$$\tilde{q}(y) = \tilde{q}(\omega^2 x) = \tilde{p}(x) = \omega[\omega^{2n} a_n x^n + \omega^{2(n-1)} a_{n-1} x^{n-1} + \dots + \omega^2 a_1 x + a_0],$$

and obtain the polynomial approximant

$$\tilde{p}(x) = \tilde{a}_n x^n + \tilde{a}_{n-1} x^{n-1} + \dots + \tilde{a}_1 x + \tilde{a}_0$$

with coefficients  $\tilde{a}_i = \omega^{2i+1} a_i \cong [1 + (2i+1)\epsilon] a_i$ .

a) Backward stability requires  $\tilde{p}(x) = p(\tilde{x})$  for some  $\tilde{x}$  close to x,  $|\tilde{x} - x|/|x| = \mathcal{O}(\epsilon)$ . When considering  $\tilde{p}(x), p(\tilde{x}) \in \mathbb{F}$ , note that  $\tilde{p}(x) = \omega p(\tilde{x})$  for  $\tilde{x} = \omega^2 x \cong (1 + 2\epsilon)x$ , to within machine precision (recall that  $fl(x) = x(1 + \epsilon)$ ), hence Horner's scheme is backward stable.

Alternatively, if  $\tilde{p}(x) = p(\tilde{x})$  is viewed as an equality between reals,  $\tilde{p}(x), p(\tilde{x}) \in \mathbb{R}$ , solving the equation

$$f(\tilde{x}) = p(\tilde{x}) - \tilde{p}(x) = p(\tilde{x}) - \omega p(\omega^2 x) = 0,$$
(1)

gives  $\tilde{x}(x; \omega, a)$ , and an estimate of  $|\tilde{x} - x|$  is needed. Let  $x_1, ..., x_n$  be the roots of p(x) = 0. The solutions to (1) arise from perturbation of the coefficient  $a_0$  of  $p, a_0 \rightarrow a_0 - \omega p(\omega^2 x)$ . This affects the  $n^{\text{th}}$  Vieta relation which for roots  $x_1, ..., x_n$  of p(x) is

$$x_1 x_2 \dots x_n = (-1^n) a_0 / a_n,$$

while for roots  $\tilde{x}_1,...,\tilde{x}_n$  of  $\,p(\tilde{x})-\omega p(\omega^2\,x)$  is

$$\tilde{x}_1 \, \tilde{x}_2 \dots \tilde{x}_n = (-1^n) [a_0 - \omega p(\omega^2 x)] / a_n = x_1 \, x_2 \dots x_n - (-1^n) \, \omega \, p(\omega^2 x) / a_n,$$

b) Since p(x) is differentiable the relative condition number is

$$\kappa = \frac{|p'(x)|}{|p(x)|} |x|$$

Consider Taylor series  $p(0) = p(x) + p'(\xi)(0-x)$ , with  $|\xi| < |x|$ , whence  $xp'(\xi) = p(x) - p(0)$ , and

$$|xp'(\xi)| = |x| |p'(\xi)| = |p(x) - p(0)| = \left|\sum_{i=0}^{n} a_i x^i - a_0\right| = \left|\sum_{i=1}^{n} a_i x^i\right| \leqslant \sum_{i=1}^{n} |a_i x^i| \leqslant \sum_{i=0}^{n} |a_i x^i|$$

Obtain

$$\kappa = \frac{|p'(x)|}{|p(x)|} |x| \leqslant \frac{\sum_{i=0}^{n} |a_i x^i|}{|\sum_{i=1}^{n} a_i x^i|},$$

to within the error of the Taylor series  $\mathcal{O}(|x|)$ .

Alternatively, use

$$\kappa_r = \lim_{\delta \to 0} \sup_{\|\delta x\| \leqslant \delta} \frac{|\delta p| |x|}{|p| |\delta x|},\tag{2}$$

and valuate

$$|\delta p| = |\tilde{p}(x) - p(x)| = \left|\sum_{i=0}^{n} (\tilde{a}_i - a_i)x^i\right| = \left|\sum_{i=0}^{n} i\epsilon a_i x^i\right| \leqslant \left|n\epsilon\sum_{i=0}^{n} a_i x^i\right| \leqslant n\epsilon\sum_{i=0}^{n} |a_i x^i|,$$

and replace in (1) to obtain

$$\kappa_r \leqslant \lim_{\delta \to 0} \sup_{\|\delta x\| \leqslant \delta} \frac{\sum_{i=0}^n |a_i x^i|}{|\sum_{i=0}^n a_i x^i|} \cdot \frac{n \epsilon |x|}{|\delta x|} \cong \frac{\sum_{i=0}^n |a_i x^i|}{|\sum_{i=0}^n a_i x^i|} n |x| = \kappa n |x|$$

c) The absolute condition number for the problem  $x \to 0$  is

$$\hat{\kappa} = \lim_{\delta \to 0} \sup_{\|\delta x\| \leq \delta} \frac{\|\delta z\|}{\|\delta x\|} = \lim_{\delta \to 0} \sup_{\|\delta x\| \leq \delta} \frac{0}{\|\delta x\|}$$

For z(x) = 0, the polynomial coefficients are  $z_i = 0 \in \mathbb{R}$  with floating point representation  $a_i = \mathrm{fl}(0) \leq \epsilon$ . The degree of z(x) is zero and worst case evaluation by Horner's scheme would return  $\tilde{z}(x) = a_n = \epsilon$ . The absolute condition number is  $\kappa = \epsilon/0 = \infty$ , hence the problem is ill-posed.

d) Since z(x) = 0 is ill-posed

$$d(p,z) = \kappa =$$

2. Write an algorithm that uses Householder reflectors to reduce a symmetric matrix  $A \in \mathbb{R}^{m \times m}$  to tridiagonal form.

Solution.